

## 1. MAIN DEFINITIONS

Main reference for this talk is [BBJ, Sections 3,4]. We start from defining some basic properties of categories which will allow us define in section 2 main four (2, 1) categories which we will study.

We will always denote by  $k$  some field,  $\text{Vect}_k$  is a category of  $k$ -vector spaces,  $\text{Vect}_{k,f.d.}$  is a category of finite dimensional vector spaces.

### Definition 1.1

A category  $\mathcal{C}$  is called  $k$ -linear if for any two objects  $X, Y \in \text{Ob}(\mathcal{C})$  a class  $\text{Hom}_{\mathcal{C}}(X, Y)$  is equipped with a  $k$ -linear structure. Such that for any  $X, Y, Z \in \text{Ob}(\mathcal{C})$  the map  $\text{Hom}(X, Y) \times \text{Hom}(Y, Z) \rightarrow \text{Hom}(X, Z)$  is  $k$ -bilinear.

Functor  $\mathcal{F}: \mathcal{C} \rightarrow \mathcal{D}$  between two  $k$ -linear categories is called  $k$ -linear if for any  $a, b \in \text{Ob}(\mathcal{C})$  the morphism  $\text{Hom}_{\mathcal{C}}(X, Y) \rightarrow \text{Hom}_{\mathcal{D}}(\mathcal{F}(X), \mathcal{F}(Y))$  is  $k$ -linear.

### Remark 1.2

One can easily show that for any  $k$ -linear (additive) category  $\mathcal{C}$  there exists a canonical bilinear functor  $\text{Vect}_{k,f.d.} \times \mathcal{C} \rightarrow \mathcal{C}$ ,  $(V, X) \mapsto V \otimes X$  such that for  $X \in \text{Ob}(\mathcal{C})$  the object  $V \otimes_k X$  represents the functor  $\mathcal{C} \rightarrow \text{Vect}_k$ ,  $Y \mapsto \text{Hom}_{\text{Vect}_k}(V, \text{Hom}_{\mathcal{C}}(X, Y))$  i.e.

$$\text{Hom}_{\mathcal{C}}(V \otimes_k X, Y) \simeq \text{Hom}_{\text{Vect}_k}(V, \text{Hom}_{\mathcal{C}}(X, Y)).$$

To construct this object we can use a basis in  $V$  and the fact that it represents some functor shows that this object is actually canonical.

### Definition 1.3

A category  $\mathcal{C}$  is called small if both objects and Hom-spaces of  $\mathcal{C}$  are sets.

### Definition 1.4

A category  $\mathcal{C}$  is called essentially small if it is equivalent to a small category.

Let us now discuss colimits. The following lemma is very useful.

### Lemma 1.5

Consider a small diagram  $J: \mathcal{I} \rightarrow \mathcal{A}$  (i.e. category  $\mathcal{I}$  is small). We denote by  $\text{Arr}(\mathcal{I})$  the set of arrows of  $\mathcal{I}$ . For an arrow  $a \in \text{Mor}(\mathcal{I})$  we denote by  $s(a), t(a) \in \text{Ob}(\mathcal{I})$  its start and target respectively. Set  $X := \coprod_{a \in \text{Mor}(\mathcal{I})} J(s(a))$ ,  $Y := \coprod_{i \in \text{Ob}(\mathcal{I})} J(i)$  (we assume that they exist). We have two morphisms  $\psi, \phi: X \rightarrow Y$  defined as follows:  $\psi|_{J(s(a))} := \iota_{t(a)} \circ a$ ,  $\phi|_{J(s(a))} := \iota_{s(a)}$ , where  $\iota_{s(a)}: J(s(a)) \rightarrow Y$ ,  $\iota_{t(a)}: J(t(a)) \rightarrow Y$  are the natural maps. Then the co-limit of  $\mathcal{F}$  along  $\mathcal{I}$  is exactly a co-equalizer of the pair  $\psi, \phi: X \rightarrow Y$  (when it exists).

### Definition 1.6

A category  $\mathcal{C}$  is called cocomplete (resp. finite cocomplete) if it contains all small (resp. finite) co-limits.

**Lemma 1.7**

Category  $\mathcal{C}$  is cocomplete (resp. finite cocomplete) iff it contains all small (resp. finite) coproducts and coequalizers.

*Proof.* Follows from lemma 1.5. □

**Example 1.8.** An example of cocomplete category is the category **Set**. One can show that if  $\mathcal{D}$  is cocomplete and  $\mathcal{C}$  is any other category then category  $[\mathcal{C}, \mathcal{D}]$  of functors from  $\mathcal{C}$  to  $\mathcal{D}$  is cocomplete (one can compute colimits pointwisely). In particular any (small) category  $\mathcal{C}$  can be fully faithfully embedded (via Yoneda) in a cocomplete category  $[\mathcal{C}^{opp}, \mathbf{Set}]$ .

**Example 1.9.** By lemma 1.5 any abelian category  $\mathcal{A}$  is finite cocomplete. Indeed it's enough to show that finite coproducts and coequalizers exist in  $\mathcal{A}$ . Existence of finite coproducts is one of the axioms of abelian category, coequalizer of two arrows  $\phi, \psi: A \rightarrow A', A, A' \in \mathcal{A}$  is nothing else but  $\text{coker}(\phi - \psi)$ .

**Example 1.10.** An example of a not cocomplete category but finite cocomplete category is the category  $A\text{-mod}_{f.g.}$  of finitely generated modules over a noetherian ring  $A$ . Being abelian it is finite cocomplete but it does not contain a colimit of the following diagram  $A \hookrightarrow A^{\oplus 2} \hookrightarrow \dots$  (which should be  $A^{\oplus \infty}$ ).

**Definition 1.11**

A non-empty category  $\mathcal{I}$  is called filtered if

- (i) for every two objects  $i, j \in \text{Ob}(\mathcal{I})$  there exists an object  $l$  and two morphisms  $i \rightarrow l, j \rightarrow l$ ,
- (ii) For every two morphisms  $u, v: i \rightarrow j$  there exists an object  $l \in \text{Ob}(\mathcal{I})$  and an arrow  $w: j \rightarrow l$  such that  $w \circ v = w \circ u$ .

**Example 1.12.** Let  $I$  be a directed set i.e. a set equipped with a preorder  $\leq$  such that any finite subset of  $I$  has an upper bound. Then we can construct a filtered category  $\mathcal{I}$  as a category whose objects are elements of  $I$  and the set  $\text{Hom}(a, b)$  consists of one element  $a \rightarrow b$  if  $a \leq b$  and is empty otherwise.

**Lemma 1.13**

Category  $\mathcal{C}$  is cocomplete iff it contains all finite and filtered coproducts.

*Proof.* Indeed, if  $\mathcal{C}$  contains all finite coproducts then it also contains coequalizers and all finite coproducts. Now any small coproduct is a filtered colimit of finite coproduct so we are done by lemma 1.5. □

**Definition 1.14**

A category  $\mathcal{C}$  is called presentable (locally-presentable) if it is cocomplete and there exists a small subset  $S$  of  $\text{Ob } \mathcal{C}$  such that any object of  $\mathcal{C}$  is a filtered colimit of objects in  $S$ .

**Example 1.15.** Examples of presentable categories include category **Set** and more generally categories  $[\mathcal{C}, \mathbf{Set}]$  where  $\mathcal{C}$  is small. Other example of presentable category is a category  $\text{Op}(X)$  for a topological space  $X$ , objects of  $\text{Op}(X)$  are open subsets

$U \subset X$  and morphisms are open embeddings, coproducts correspond to unions and coequalizers are trivial. It is an exercise to show that the category of coalgebras over a field  $k$  is presentable.

**Definition 1.16**

A functor  $\mathcal{F}: \mathcal{A} \rightarrow \mathcal{B}$  from a cocomplete category  $\mathcal{A}$  is called *cocontinuous* if it preserves colimits.

**Remark 1.17**

Note that in the same way as in the proof of lemma 1.13 using lemma 1.5 we see that a functor is cocontinuous iff it preserves finite and filtered colimits.

**Example 1.18.** Let  $X, Y$  be two topological spaces and  $\text{Op}(X), \text{Op}(Y)$  are the corresponding categories of open subsets. Then any continuous map  $f: Y \rightarrow X$  defines a cocontinuous functor  $f^*: \text{Op}(X) \rightarrow \text{Op}(Y)$  by sending  $U \in \text{Op}(X)$  to  $f^{-1}(U) \in \text{Op}(Y)$ .

**Definition 1.19**

An object  $C \in \mathcal{C}$  of a category  $\mathcal{C}$  which admits all filtered colimits is called *compact* if the functor  $\text{Hom}_{\mathcal{C}}(C, \bullet)$  commutes with filtered colimits. We denote by  $\mathcal{C}_c \subset \mathcal{C}$  the full subcategory consisting of compact objects of  $\mathcal{C}$ .

**Example 1.20.** Compact objects in the category **Set** are precisely finite sets.

**Example 1.21.** Let  $R$  be a noetherian ring and  $\mathcal{C} = R\text{-mod}$  the category of  $R$ -modules. Then an object  $C \in \mathcal{C}$  is compact iff it is finitely generated.

*Proof.* The implication  $\Leftarrow$  is an exercise. Let us prove the implication  $\Rightarrow$ . We fix a compact object  $M \in \mathcal{C}$ . Note that we can present  $M$  as a colimit of its finitely generated submodules  $M_i \subset M$ . We have  $\text{Hom}(M, M) = \text{Hom}(M, \text{colim } M_i) = \text{colim}(\text{Hom}(M, M_i))$ . Consider now the element  $id \in \text{Hom}(M, M)$ . We see that there exists  $M_i \subset M$  and  $f: M \rightarrow M_i$  such that  $id = \iota_i \circ f$ , where  $\iota_i: M_i \hookrightarrow M$  is the embedding. It follows that  $M = M_i$ , hence,  $M$  is finitely presented.  $\square$

**Example 1.22.** Let  $X$  be a topological space and recall a category  $\text{Op}(X)$ . Then an object  $C \in \text{Op}(X)$  is compact iff it is compact as a topological space.

*Proof.* Fix a compact object  $C \in \text{Op}(X)$ . Consider a covering  $C = \bigsqcup_{i \in I} U_i$  by open subsets. For any finite subset  $K \subset I$  define  $U_K := \bigcup_{i \in K} U_i$ . The set  $\{U_K\}$  together with natural open embeddings defines a filtered system. Note that  $\text{colim}_K U_K = C$ . We see that  $\text{Hom}(C, C) = \text{Hom}(C, \text{colim}_K U_K) = \text{colim}_K \text{Hom}(C, U_K)$ . It follows that  $C = U_K$  for some  $K$  i.e.  $C$  is compact.

It is an exercise to check that any compact subspace  $C \subset X$  is compact in  $\text{Op}(X)$ .  $\square$

**Remark 1.23**

Let us point out that restricting ourselves to filtered colimits is crucial in the definition of compact object. For example in the category  $A\text{-mod}$  compact objects are precisely finitely presented modules while objects  $P \in A\text{-mod}$  such that  $\text{Hom}_{\mathcal{C}}(C, \bullet)$  commutes with all colimits are projective finitely presented modules.

**Remark 1.24**

Any functor  $\mathcal{F}: \mathcal{A} \rightarrow \mathcal{B}$  commutes with finite filtered colimits. Indeed if  $\mathcal{I}$  is a finite filtered category and  $J: \mathcal{I} \rightarrow \mathcal{A}$  is a diagram then there is always the maximal element  $i \in \text{Ob}(\mathcal{I})$ . Now it follows from the definitions that  $\text{colim } J = J(i)$  i.e.  $\mathcal{F}(\text{colim } J) = \mathcal{F}(J(i)) = \text{colim } \mathcal{F} \circ J$ .

**Lemma 1.25**

Let  $\mathcal{C}$  be a cocomplete category then a finite colimit of compact objects is compact.

*Proof.* Indeed assume that  $C = \text{colim}_i C_i$ , with  $C_i$ -compact and such that the indexing set is finite. Consider now arbitrary filtered colimit  $\text{colim}_j X_j$ . We have

$$\begin{aligned} \text{Hom}(C, \text{colim}_j X_j) &= \text{Hom}(\text{colim}_i C_i, \text{colim}_j X_j) = \lim_i \text{Hom}(C_i, \text{colim}_j X_j) = \\ &= \lim_i \text{colim}_j \text{Hom}(C_i, X_j) = \text{colim}_j \lim_i \text{Hom}(C_i, X_j) = \\ &= \text{colim}_j \text{Hom}(\lim_i C_i, X_j) = \text{colim}_j \text{Hom}(C, X_j), \end{aligned}$$

here we use the following standard fact – finite limits commute with filtered colimits (exercise).  $\square$

**Definition 1.26**

A functor  $\mathcal{F}: \mathcal{A} \rightarrow \mathcal{B}$  is called compact if it sends compact objects to compact objects.

**Example 1.27.** A continuous map  $f: Y \rightarrow X$  induces a compact functor  $f^*: \text{Ob}(X) \rightarrow \text{Ob}(Y)$  iff  $f$  is proper.

**Definition 1.28**

An object  $S \in \text{Ob}(\mathcal{C})$  is called a generator if for every pair of morphisms  $f, g: X \rightarrow Y$  in  $\mathcal{C}$ , if  $f \circ l = g \circ l$  for every morphism  $l: S \rightarrow X$  then  $f = g$ .

**Example 1.29.** Let  $\mathcal{C}$  be a category of  $A$ -modules. Then  $A \in \text{Ob}(\mathcal{C})$  is a generator. Indeed if  $M, N$  are two  $A$ -modules and  $f, g: M \rightarrow N$  are two morphisms then to any  $m \in M$  we can associate a morphism  $l_m: A \rightarrow M$ ,  $a \mapsto am$  then from  $f \circ l_m = g \circ l_m$  we deduce  $f(m) = g(m)$ .

**Definition 1.30**

A Grothendieck category is an abelian cocomplete category which has a generator and such that filtered colimits are exact.

**Example 1.31.** Let  $A$  be a  $k$ -algebra then the category  $A\text{-mod}$  is Grothendieck. More generally if  $(X, \mathcal{O}_X)$  is a ringed space then the category of sheaves of  $\mathcal{O}_X$ -modules is Grothendieck.

Let us now formulate adjoint functor theorem in the setting of presentable categories (see [AR]).

**Proposition 1.32**

Let  $\mathcal{F}: \mathcal{A} \rightarrow \mathcal{B}$  be a functor between locally presentable categories  $\mathcal{A}, \mathcal{B}$ . Then  $\mathcal{F}$  admits a right adjoint iff it preserves all small co-limits.

**Definition 1.33**

Let  $\mathcal{F}: \mathcal{A} \rightarrow \mathcal{B}$  be a functor. Suppose that all finite limits (resp. co-limits) exist in  $\mathcal{A}$ . We say that  $\mathcal{F}$  is left (resp. right) exact if it commutes with finite limits (resp. colimits).

**Lemma 1.34**

Functor  $\mathcal{F}: \mathcal{A} \rightarrow \mathcal{B}$  is right (resp. left) exact iff it commutes with finite products (resp. coproducts) and equalizers (resp. coequalizers).

*Proof.* The implication  $\Rightarrow$  is clear. The opposite implication follows from lemma 1.5.  $\square$

**Corollary 1.35**

For abelian categories  $\mathcal{A}, \mathcal{B}$  the two notions of right (resp. left) exact functors coincide.

## 2. MAIN CATEGORIES

Let  $\mathcal{V}$  be a monoidal category. A  $\mathcal{V}$ -category  $\mathcal{A}$  (or a category enriched over  $\mathcal{V}$ ) is

- (i) a class of objects  $\text{Ob}(\mathcal{A})$ ,
- (ii) for any  $X, Y \in \text{Ob}(\mathcal{A})$  an object  $\text{Hom}(X, Y) \in \mathcal{V}$ ,
- (iii) for each  $X, Y, Z \in \text{Ob}(\mathcal{A})$  a morphism

$$\circ_{XYZ}: \text{Hom}(X, Y) \otimes \text{Hom}(Y, Z) \rightarrow \text{Hom}(X, Z),$$

- (iv) for each  $X \in \text{Ob}(\mathcal{A})$  a morphism  $id_X: 1 \rightarrow \text{Hom}(X, X)$

such that the following diagrams commute:

$$\begin{array}{ccc} (\text{Hom}(Z, H) \otimes \text{Hom}(Y, Z)) \otimes \text{Hom}(X, Y) & \xrightarrow{\circ_{YZH} \otimes id} & \text{Hom}(Y, H) \otimes \text{Hom}(X, Y), \\ \downarrow \alpha & & \downarrow \circ_{XYH} \\ \text{Hom}(Z, H) \otimes (\text{Hom}(Y, Z) \otimes \text{Hom}(X, Y)) & \xrightarrow{id \otimes \circ_{XYZ}} & \text{Hom}(Z, H) \otimes \text{Hom}(X, Z) \\ & & \uparrow \circ_{XZH} \end{array}$$

$$\begin{array}{ccc} 1 \otimes \text{Hom}(X, Y) & \xrightarrow{id_Y \otimes id} & \text{Hom}(Y, Y) \otimes \text{Hom}(X, Y), \\ & \searrow \lambda & \swarrow \circ_{XYX} \\ & \text{Hom}(X, Y) & \\ \\ \text{Hom}(X, Y) \otimes 1 & \xrightarrow{id \otimes id_X} & \text{Hom}(X, Y) \otimes \text{Hom}(X, X), \\ & \searrow \rho & \swarrow \circ_{XXY} \\ & \text{Hom}(X, Y) & \end{array}$$

where  $\alpha, \lambda, \rho$  are the natural morphisms in the tensor category  $\mathcal{V}$ .

**Example 2.36.** A category enriched over  $\mathbf{Vect}_k$  is nothing else but  $k$ -linear category. A category enriched over  $\mathbf{Set}$  is nothing else but a small category.

Recall that if we have two categories  $\mathcal{C}, \mathcal{D}$  then we can form their cartesian product whose objects are  $\text{Ob}(\mathcal{C}) \times \text{Ob}(\mathcal{D})$  and  $\text{Hom}_{\mathcal{C} \times \mathcal{D}}((C, D), (C', D')) := \text{Hom}_{\mathcal{C}}(C, C') \times \text{Hom}_{\mathcal{D}}(D, D')$  for  $C, C' \in \text{Ob}(\mathcal{C}), D, D' \in \text{Ob}(\mathcal{D})$ .

**Definition 2.37**

A (strict) 2-category is a category enriched over  $\mathbf{Cat}$  with monoidal structure given by cartesian product.

**Definition 2.38**

A  $(2, 1)$ -category is a 2-category in which any 2-morphism is invertible.

Our main players will be the following four  $(2, 1)$ -categories:

(1) **Rex** is a category of essentially small finitely cocomplete  $k$ -linear categories with morphisms right exact functors and 2-morphisms –  $k$ -linear natural isomorphisms.

(2) **Pr** is a category of presentable  $k$ -linear categories with morphisms cocontinuous functors and 2-morphisms –  $k$ -linear natural isomorphisms.

(3) **Pr<sub>c</sub>** is a category of  $k$ -linear cocomplete categories  $\mathcal{C}$  such that  $\mathcal{C}_c$  is essentially small and any object of  $\mathcal{C}$  is a filtered limit of compact objects. Morphisms between such categories are compact cocontinuous functors and 2-morphisms –  $k$ -linear natural isomorphisms.

(4) **Gr** is a category of Grothendieck categories.

**Remark 2.39**

Note that by proposition 1.32 morphisms in categories **Pr**, **Pr<sub>c</sub>** have right adjoints.

**Example 2.40.** One very important example of an object of **Rex** is a category  $C - \text{comod}_{f.d.}$  of finite dimensional comodules over some  $k$ -coalgebra  $C$ .

Now our goal is to construct an equivalence  $\mathbf{Rex} \simeq \mathbf{Pr}_c$ . Starting from a small finitely co-complete category  $\mathcal{C}$  we can construct its ind-completion as follows. Its objects are functors  $\mathcal{F}: \mathcal{I} \rightarrow \mathcal{C}$  from a small filtered category  $\mathcal{I}$  to  $\mathcal{C}$ . Morphisms in  $\text{ind}(\mathcal{C})$  are natural transformations of functors. We have a canonical fully faithful embedding  $\iota: \mathcal{C} \hookrightarrow \text{ind}(\mathcal{C})$ . The following lemma (exercise) describes a universal property of the category  $\text{ind}(\mathcal{C})$ .

**Lemma 2.41**

For any category  $\mathcal{D}$  which has all filtered colimits and a functor  $\mathcal{F}: \mathcal{C} \rightarrow \mathcal{D}$  there exists a unique functor  $\text{ind}(\mathcal{F}): \text{ind}(\mathcal{C}) \rightarrow \mathcal{D}$  which preserves filtered colimits and such that  $\mathcal{F} = \text{ind}(\mathcal{F}) \circ \iota$ .

The following is true.

**Lemma 2.42**

(a) For any  $C \in \text{Ob}(\mathcal{C})$  the object  $\iota(C) \in \text{ind}(\mathcal{C})$  is compact.

(b) If  $\tilde{C} \in \text{ind}(\mathcal{C})$  is compact then  $\tilde{C}$  is a retract of some object  $C$  of  $\mathcal{C}$  i.e. there exist morphisms  $\tilde{C} \xrightarrow{\iota} C \xrightarrow{\pi} \tilde{C}$  such that  $\pi \circ \iota = \text{id}_{\tilde{C}}$ .

(c) If  $\mathcal{C}$  is finite cocomplete then any compact object of  $\tilde{\mathcal{C}}$  is isomorphic to  $\iota(C)$  for some  $C \in \text{Ob}(\mathcal{C})$ .

*Proof.* Let us prove (a). Fix an object  $C \in \mathcal{C}$  and let  $\tilde{J}: \mathcal{I} \rightarrow \mathcal{C}, i \mapsto X_i$  be a filtered diagram. From the definitions it follows that it is enough to deal with the diagrams  $\tilde{J} = \iota \circ J, J: \mathcal{I} \rightarrow \mathcal{C}$ . We have to show that  $\text{Hom}_{\tilde{\mathcal{C}}}(\iota(C), \text{colim } X_i) = \text{colim}(\text{Hom}(C, X_i))$ . So we need to check a universal property of colimit for  $\text{Hom}_{\tilde{\mathcal{C}}}(\iota(C), \text{colim } X_i)$ . Take a set  $Z$  together with compatible homomorphisms  $\psi_i: \text{Hom}(C, X_i) \rightarrow Z$ . We need to construct a morphism  $\text{Hom}(C, \text{colim } X_i) \rightarrow Z$ . Note that one can consider  $\text{colim } X_i$  as an object of  $\text{ind}(\mathcal{C})$ . By the definition an element of  $\text{Hom}(C, \text{colim } X_i)$  is a family of compatible morphisms  $\varphi_i: C \rightarrow X_i$ . We now construct a morphism  $\text{Hom}(C, \text{colim } X_i) \rightarrow Z$  by sending  $(\varphi_i)$  to an element  $\psi_i(\varphi_i)$  (it does not depend on  $i$ ). The claim follows.

Let us now prove part (b). Fix a compact object  $\tilde{C} \in \text{ind}(\mathcal{C})$ . We can write  $\tilde{C} = \text{colim}_{i \in \text{Ob}(\mathcal{I})}(X_i)$  for some diagram  $\mathcal{I} \rightarrow \mathcal{C}, i \mapsto X_i$ . We see that  $\text{Hom}_{\text{ind}(\mathcal{C})}(\tilde{C}, \tilde{C}) = \text{Hom}_{\text{ind}(\mathcal{C})}(\tilde{C}, \text{colim}_{i \in \text{Ob}(\mathcal{I})}(X_i)) = \text{colim}_i(\tilde{C}, X_i)$  so, in particular, morphism  $\text{id}: \tilde{C} \rightarrow \tilde{C}$  can be decomposed as  $\text{id} = \pi \circ \iota$  for some  $X_i \in \mathcal{C}$ .

To prove (c) it remains to note that  $\tilde{C}$  from (b) is isomorphic to a coequalizer of the pair  $\iota \circ \pi, \text{id}: C \rightarrow C$  (we keep notations from the proof of (b)) which lies in  $\mathcal{C}$  since it is finite cocomplete.  $\square$

### Proposition 2.43

We have an equivalence of  $(2, 1)$ -categories  $\mathbf{Rex} \simeq \mathbf{Pr}_c$ .

*Proof.* Fix a category  $\mathcal{C} \in \mathbf{Rex}$ . We can consider its ind-completion  $\text{ind}(\mathcal{C})$ . We claim that  $\text{ind}(\mathcal{C}) \in \mathbf{Pr}_c$ . Indeed category  $\text{ind}(\mathcal{C})$  contains all finite (because  $\mathcal{C}$  was finite cocomplete and colimits commute with each other) and filtered colimits so by lemma 1.13 it is cocomplete, note also that by lemma 2.42 any object  $X \in \mathcal{C}$  is compact as an object of  $\text{ind}(\mathcal{C})$ . Note also that if  $\mathcal{F}: \mathcal{C} \rightarrow \mathcal{C}'$  is a morphism in the category  $\mathbf{Rex}$  i.e. a functor which commutes with finite colimits then  $\text{ind}(\mathcal{F})$  is cocontinuous.

So we obtain a functor  $\text{ind}: \mathbf{Rex} \rightarrow \mathbf{Pr}_c$ .

Let us construct an inverse functor  $\mathbf{Pr}_c \rightarrow \mathbf{Rex}$ . It sends  $\mathcal{D} \in \mathbf{Pr}_c$  to a full subcategory  $\mathcal{D}_c \subset \mathcal{D}$  of compact objects. Recall that by the definition  $\mathcal{D}_c$  is essentially small. Note also that by lemma 1.25 the category  $\mathcal{D}_c$  is finite cocomplete. It follows that  $\mathcal{D}_c \in \mathbf{Rex}$ . We obtain a functor  $\text{comp}: \mathbf{Pr}_c \rightarrow \mathbf{Rex}_c$ .

It follows from lemma 2.42 that  $\text{ind}, \text{comp}$  are mutually inverse equivalences.  $\square$

We are now going to investigate objects  $\mathcal{C}, \mathcal{D} \in \mathbf{Pr}_c \simeq \mathbf{Rex}$  and functors  $L: \mathcal{C} \rightarrow \mathcal{D}$ . For that we recall that by proposition 1.32 functor  $L$  admits a right adjoint  $R: \mathcal{D} \rightarrow \mathcal{C}$ . We will investigate category  $\mathcal{D}$  using the functor  $T := R \circ L: \mathcal{C} \rightarrow \mathcal{D}$  (it will have a structure of a monad) and modules over  $T$ . Let us start from the general definition.

### Definition 2.44

Let  $(\mathcal{C}, \otimes)$  be a monoidal category. An object  $A \in \mathcal{C}$  is called a unital algebra object if we have morphisms  $m: A \otimes A \rightarrow A, i: 1 \rightarrow A$  such that standard axioms of associative

unital algebra holds:

$$\begin{array}{ccccc}
 (A \otimes A) \otimes A & \xrightarrow{\alpha} & A \otimes (A \otimes A) & \xrightarrow{id \otimes m} & A \otimes A, \\
 & \searrow m \otimes id & & & \swarrow m \\
 & & A \otimes A & \xrightarrow{m} & A
 \end{array}$$
  

$$\begin{array}{ccccc}
 1 \otimes A & \xrightarrow{i \otimes id} & A \otimes A & \xleftarrow{id \otimes i} & A \otimes A. \\
 & \searrow \lambda & \downarrow m & \swarrow \rho & \\
 & & A & & 
 \end{array}$$

**Example 2.45.** Let  $\mathcal{C} = \text{Vect}_k$  be a category of vector spaces over  $k$ . Then a unital algebra object of  $\mathcal{C}$  is exactly an associative unital  $k$ -algebra.

To each  $k$ -linear category  $\mathcal{C}$  we can associate the tensor category  $(\text{End}(\mathcal{C}), \circ)$  whose objects are functors  $\mathcal{F}: \mathcal{C} \rightarrow \mathcal{C}$ , morphisms are natural transformations and tensor structure comes from the composition of morphisms.

**Definition 2.46**

A monad  $T$  in a category  $\mathcal{C}$  is a unital algebra object in the category  $\text{End}(\mathcal{C})$ .

**Example 2.47.** Let  $L: \mathcal{C} \rightarrow \mathcal{D}$ ,  $R: \mathcal{D} \rightarrow \mathcal{C}$  be an adjoint pair of functors. Then the composition  $T = R \circ L$  is a monad on  $\mathcal{C}$  via the adjunction and counit

$$\eta: id_{\mathcal{C}} \rightarrow R \circ L, \epsilon: L \circ R \rightarrow id_{\mathcal{D}}.$$

**Definition 2.48**

If  $T$  is a monad in the category  $\mathcal{C}$  with multiplication  $m: T^2 \rightarrow T$  and unit  $\eta: id_{\mathcal{C}} \rightarrow T$  then we define the category  $T\text{-mod}_{\mathcal{C}}$  as a category of pairs  $(X, f)$ , where  $X \in \mathcal{C}$  and  $f$  is a morphism  $T(X) \xrightarrow{f} X$  such that the following diagrams commute:

$$\begin{array}{ccc}
 T^2(X) & \xrightarrow{m(X)} & T(X), \quad X \xrightarrow{\eta(x)} T(X). \\
 \downarrow T(f) & & \downarrow f \\
 T(X) & \xrightarrow{f} & X
 \end{array}$$

Morphisms between  $(X, f), (X', f') \in \text{Ob}(T - \text{mod}_{\mathcal{C}})$  are  $h: X \rightarrow X'$  such that  $h \circ f = f' \circ T(h)$ .

**Example 2.49.** Functor  $id_{\mathcal{C}}: \mathcal{C} \rightarrow \mathcal{C}$  is a monad via the natural isomorphisms  $id_{\mathcal{C}} \xrightarrow{\sim} id_{\mathcal{C}}, id_{\mathcal{C}} \circ id_{\mathcal{C}} \xrightarrow{\sim} id_{\mathcal{C}}$ . We have an equivalence  $id_{\mathcal{C}} - \text{mod}_{\mathcal{C}} \xrightarrow{\sim} \mathcal{C}$  given by  $X \mapsto (X, id_{\mathcal{C}}(X))$ .

**Example 2.50.** More generally assume that  $L, R$  are mutually inverse equivalences. Then the morphisms  $\eta: id_{\mathcal{C}} \rightarrow R \circ L, \epsilon: L \circ R \rightarrow id_{\mathcal{D}}$  are isomorphisms and we see that  $T - \text{mod}_{\mathcal{C}}$  is equivalent to  $\mathcal{C}$  via a morphism  $\mathcal{C} \rightarrow T - \text{mod}_{\mathcal{C}}, X \mapsto (X, \eta^{-1}(X))$ , the inverse functor is a forgetting functor  $\text{forg}: T - \text{mod}_{\mathcal{C}} \rightarrow \mathcal{C}$ .



We always have a functor  $\tilde{R}: \mathcal{D} \rightarrow T\text{-mod}_{\mathcal{C}}$  sending  $X \in \mathcal{D}$  to  $R(X) \in \mathcal{C}$  equipped with the canonical action

$$act: R \circ L \circ R(A) \xrightarrow{id \otimes \epsilon} R(A).$$

Recall that by proposition 1.32 if  $\mathcal{C}, \mathcal{D} \in \mathbf{Pr}$  that ANY morphism  $L: \mathcal{C} \rightarrow \mathcal{D}$  of  $\mathbf{Pr}$  admits a right adjoint  $R: \mathcal{D} \rightarrow \mathcal{C}$  so we obtain a functor  $\tilde{R}: \mathcal{D} \rightarrow T\text{-mod}_{\mathcal{C}}$ . Our next goal is to formulate necessary and sufficient for this functor to be an equivalence.

### 3. BARR-BECK

#### Definition 3.51

Abelian category  $\mathcal{D}$  is called *locally finite* if it has finite dimensional spaces of morphisms and each object has finite length.

#### Definition 3.52

Abelian category  $\mathcal{D}$  is called *finite* if it is locally finite has finite number of simple objects, and has enough projectives.

Let us start from the following proposition.

#### Proposition 3.53

Let  $\mathcal{D}$  be a finite abelian category. Then  $\mathcal{D} \simeq A\text{-mod}_{f.d.}$  for some finite dimensional algebra  $A$ . More precisely  $A = \text{End}(P)^{\text{opp}}$ , where  $P$  is a projective generator of  $\mathcal{D}$ .

*Proof.* Let  $P$  be a projective generator of  $\mathcal{D}$ . We have a functor  $R = \text{Hom}(P, \bullet): \mathcal{D} \rightarrow \mathcal{C}$ , where  $\mathcal{C} := \text{End}(P)^{\text{opp}}\text{-mod}_{f.d.}$ . Note that  $R$  is left exact. Let us prove that  $R$  has a left adjoint functor (it actually already follows from a version proposition 1.32). Note that if we already know that  $\mathcal{C} \simeq A\text{-mod}$  for some  $A$  than  $P$  must be an  $A$ -module and the functor  $L$  is given by  $M \mapsto P \otimes_{\text{End}(P)^{\text{opp}}} M$  for  $M \in \mathcal{C}$ .

We now just mimick the construction of tensor product. Fix a basis  $\{f_j\}$  of  $\text{End}(P)^{\text{opp}}$ , fix also a basis  $e_1, \dots, e_m$  of  $M$ . We consider an object  $P^{\oplus m}$  and want to quotient it by the elements of the form  $p \otimes f_j e_i - f_j(p) \otimes e_i$ . For that consider morphisms  $\phi_{ij}, \psi_{ij}: P \rightarrow P^{\oplus m}$ :  $\phi_{ij}$  is the map  $f_j$  composed with the  $i$ -th coordinate embedding  $P \hookrightarrow P^{\oplus m}$ , to define  $\psi_{ij}$  we decompose  $f_j e_i = \sum a_l e_l$  and define  $\psi_{ij}$  as  $a_l \text{id}_P$  on  $l$ -th summand. Then we set

$$L(M) = P \otimes_{\text{End}(P)^{\text{opp}}} M := P^{\oplus m} / \text{Span}_k(\text{Im}(\psi_{ij} - \phi_{ij} \mid i, j)).$$

Let us also describe a more canonical way of constructing a functor  $L$ . Note that we have a forgetfull functor  $\mathcal{C} \rightarrow \text{Vect}_k$  and we denote by  $M$  the corresponding vector space. Consider an object  $M \otimes_k P \in \mathcal{D}$  (see remark 1.2). We have two morphisms  $\phi, \psi: M \otimes \text{End}(P) \otimes_k P \rightarrow M \otimes_k P$  given as follows:  $\phi = \text{id}_M \otimes \text{act}_P, \psi = \text{act}_M \otimes \text{id}_P$ , where  $\text{act}_M: M \otimes_k \text{End}(P) \rightarrow M, (a, f) \mapsto af, \text{act}_P: \text{End}(P) \otimes_k P \rightarrow P, f \otimes p \mapsto f(p)$  (more formally we use the identification  $\text{Hom}(\text{End}(P) \otimes_k P, P) \simeq \text{Hom}(\text{End}(P), \text{Hom}(P, P))$  and then morphism  $\text{act}_P$  just corresponds to  $\text{id}_P$ ). We then can define  $L(M) := \text{coker}(\phi - \psi)$ .

It is easy to see that  $L$  is left adjoint to  $R$ .

It remains to check that the adjunction morphisms  $\eta: \text{id}_{\mathcal{C}} \rightarrow R \circ L$ ,  $\epsilon: L \circ R \rightarrow \text{id}_{\mathcal{D}}$  are isomorphisms. This can be done on the generators  $P, \text{Hom}(P, P)$  of the categories  $\mathcal{D}, \mathcal{C}$  respectively.

We have

$$\begin{aligned} L(R(P)) &= L(\text{Hom}(P, P)) = P \otimes_{\text{End}(P)^{\text{opp}}} \text{Hom}(P, P) = P, \\ R(L(\text{Hom}(P, P))) &= R(P) = \text{Hom}(P, P). \end{aligned}$$

The claim follows.  $\square$

**Example 3.54.** Let us give couple examples of categories which satisfy the conditions of proposition 3.53. One important class of examples are blocks of BGG-category  $\mathcal{O}$  for semi-simple finite dimensional Lie algebra  $\mathfrak{g}$ . For  $\mathfrak{g} = \mathfrak{sl}_2$  one can easily describe the corresponding algebra  $\text{End}(P)^{\text{opp}}$  explicitly.

The following generalization of proposition 3.53 to the case of locally finite abelian categories holds (the proof is not very hard, see [EGNO, Section 1.10]).

**Proposition 3.55**

*Let  $\mathcal{D}$  be a locally finite abelian category. Then  $\mathcal{D} \simeq D - \text{comod}_{f,d}$  where the later is a category of finite dimensional comodules over a unique coalgebra  $D$ . If  $\mathcal{D}$  is finite then  $D$  is finite dimensional and  $\mathcal{D} \simeq D^* - \text{mod}_{f,d}$ . (c.f. proposition 3.53).*

**Remark 3.56**

Note that any locally finite category lies in **Rex**. The subclass of locally finite categories in **Rex** is very important, for example, it is closed under the Deligne-Kelly tensor product  $\boxtimes$  on **Rex** (see section 4 and remark 4.67).

Let us now generalize proposition 3.53 to our setting.

**Definition 3.57**

*A fork in a category  $\mathcal{A}$  is a triple  $f, g: A \rightarrow B$ ,  $e: B \rightarrow C$  such that  $ef = eg$ . A split coequalizer is a fork together with morphisms  $t: B \rightarrow A$ ,  $s: C \rightarrow B$  such that  $es = \text{id}_C$ ,  $se = gt$ ,  $ft = \text{id}_B$ .*

The following lemma is easy.

**Lemma 3.58**

*Let  $(f, g, e)$  be a split coequalizer. Then for any functor  $\mathcal{F}: \mathcal{A} \rightarrow \mathcal{B}$  the morphism  $\mathcal{F}(e)$  is a coequalizer of the pair  $\mathcal{F}(f), \mathcal{F}(g)$ .*

*Proof.* Note that for any functor  $\mathcal{F}$ ,  $(\mathcal{F}(f), \mathcal{F}(g), \mathcal{F}(e))$  is a split coequalizer. So it remains to check that if  $(f, g, e)$  is a split coequalizer then  $e$  is a coequalizer of  $f, g$ . Indeed if  $h: B \rightarrow D$  is some morphism such that  $h \circ f = h \circ g$  then  $h \circ s: C \rightarrow D$  gives us the desired morphism. If  $p: C \rightarrow D$  is any other map which makes diagram commutative then we must have  $p \circ e = h$  so  $p = p \circ e \circ s = h \circ s$  and uniqueness follows.  $\square$

The following theorem is a generalization of the proposition 3.53. Recall adjoint functors  $L: \mathcal{C} \rightarrow \mathcal{D}$ ,  $R: \mathcal{D} \rightarrow \mathcal{C}$ .

**Theorem 3.59**

The functor  $\tilde{R}: \mathcal{D} \rightarrow T - \text{mod}_{\mathcal{C}}$  is an equivalence if and only if

- (1)  $R$  is conservative i.e. if  $f: X \rightarrow Y$  in  $\mathcal{D}$  is such that  $R(f)$  is an isomorphism then  $f$  is an isomorphism,
- (2)  $\mathcal{D}$  has coequalizers of  $R$ -split parallel pairs (those parallel pairs of morphisms in  $\mathcal{D}$ , which  $R$  sends to pairs having a split coequalizer in  $\mathcal{D}$ ) and  $R$  preserves those coequalizers.

*Proof.* Let us prove the implication  $\Rightarrow$ . We assume that  $\mathcal{D}$  is equivalent to  $T - \text{mod}_{\mathcal{C}}$  via  $\tilde{R}$ . After this equivalence the functor  $R: \mathcal{D} \rightarrow \mathcal{C}$  becomes isomorphic to a forgetfull functor  $\text{Forg}: T - \text{mod}_{\mathcal{C}} \rightarrow \mathcal{C}$  which is obviously conservative.

Let us now fix two maps  $h, l: (X, f) \rightarrow (Y, g)$  in  $T - \text{mod}_{\mathcal{C}}$  such that the corresponding maps  $h, l: X \rightarrow Y$  have a split coequalizer  $e: Y \rightarrow Z$  in  $\mathcal{C}$ . We need to construct a coequalizer for  $h, l: (X, f) \rightarrow (Y, g)$ . We have the following diagram

$$\begin{array}{ccccc} T(X) & \xrightarrow{T(h), T(l)} & T(Y) & \xrightarrow{T(e)} & T(Z) \\ \downarrow f & & \downarrow g & & \\ X & \xrightarrow{h, l} & Y & \xrightarrow{e} & Z \end{array}$$

Note that by lemma 3.58,  $T(e)$  is a coequalizer. It follows from the definitions that  $egT(l) = elf = ehf = egT(h)$ . We conclude that there exists  $m: T(Z) \rightarrow Z$  such that the diagram commute. It is an exercise to check that  $m$  is compatible with a monad (algebra) structure. It remains to show that  $e: (Y, g) \rightarrow (Z, m)$  is a coequalizer.

Consider any map  $d: (Y, g) \rightarrow (Q, p)$  in  $T - \text{mod}_{\mathcal{C}}$  such that  $dl = dh$ . Recall that  $e: Y \rightarrow Z$  is a coequalizer in  $\mathcal{C}$  so there exists a unique  $d': Z \rightarrow Q$  such that  $d'e = d$ . Using the fact that  $T(e)$  is a coequalizer (uniqueness part) we obtain  $p \circ T(d') = d' \circ m$  (maps  $p \circ T(d'), d' \circ m: T(Z) \rightarrow Q$  satisfy  $p \circ T(d') \circ T(h) = d' \circ m \circ T(h), p \circ T(d') \circ T(l) = d' \circ m \circ T(l)$  so they must coincide). So we get a desired map  $d': (Z, m) \rightarrow (Q, p)$ .

Let us prove the implication  $\Leftarrow$ . Recall unit and counit morphisms

$$\eta: id_{\mathcal{C}} \rightarrow R \circ L, \epsilon: L \circ R \rightarrow id_{\mathcal{D}}.$$

Recall that we have a multiplication  $m: T^2 \rightarrow T$ . Let us first note that for any  $(x, f) \in \text{Ob}(T - \text{mod}_{\mathcal{C}})$  the fork  $(m(x), T(f), f)$  splits by the pair  $s = \eta(x), t = \eta(Tx)$ .

We have the forgetfull functor  $\text{Forg}_{\mathcal{C}}: T - \text{mod}_{\mathcal{C}} \rightarrow \mathcal{C}$ . We also have a functor  $\tilde{L}: \mathcal{C} \rightarrow T - \text{mod}_{\mathcal{C}}$  given by  $X \mapsto (TX, m(X))$  on the level of objects. It is easy to see that  $\tilde{L}$  is left adjoint to  $\text{Forg}_{\mathcal{C}}$ .

So we have two pairs of adjoint functors

$$(R, L), (\text{Forg}_{\mathcal{C}}, \tilde{L}), R: \mathcal{D} \rightarrow \mathcal{C}, \text{Forg}_{\mathcal{C}}: T - \text{mod}_{\mathcal{C}} \rightarrow \mathcal{C}.$$

Note that  $R \circ L = \text{Forg}_{\mathcal{C}} \circ \tilde{L}: \mathcal{C} \rightarrow \mathcal{C}$ . so these adjoint pairs define the same monads.

Let us now prove the following general lemma which will allow us to...

**Lemma 3.60**

Let  $(R, L), (R', L')$  be two adjoint pairs,  $R: \mathcal{D} \rightarrow \mathcal{C}, R': \mathcal{D}' \rightarrow \mathcal{C}$  such that  $R \circ L = R' \circ L'$ . Assume also that the condition (2) of theorem 3.59 holds for  $R$ . Then there

exists a unique functor  $Q: \mathcal{D}' \rightarrow \mathcal{D}$  such that  $RQ = R', QL' = L$  (in particular,  $Q\epsilon' = \epsilon Q$ ).

*Proof.* Let us start from the uniqueness of  $Q$ . Note that for any  $x \in \text{Ob}(\mathcal{C})$  the triple  $(LR(\epsilon(x)), \epsilon(LR(x)), \epsilon(x)), LR(\epsilon(x)), \epsilon(LR(x)): LRLR(x) \rightarrow LR(x), \epsilon(x): LR(x) \rightarrow x$ , is folk.

Let us consider a folk which corresponds to  $x = Qy$  for some  $y \in \text{Ob}(\mathcal{D}')$ . We obtain the folk  $(LR(\epsilon(Qy)), \epsilon(LR(Qy)), \epsilon(Qy)) = (LR'(\epsilon'(y)), \epsilon(LR'(y)), Q\epsilon'(y))$ .

Let us now apply  $R$  to this folk. We obtain a folk  $(RLR'(\epsilon'(y)), R\epsilon(LR'(y)), RQ\epsilon'(y))$  which splits because it coincides with the folk  $(m(R'y), T(f), f)$  for  $f = \epsilon'(R'y)$ . It follows that  $Qy$  should be a coequalizer of  $(LR'(\epsilon'(y)), \epsilon(LR'(y)))$  i.e. it is uniquely defined. The uniqueness of  $Q$  follows.

Let us prove the existence of  $Q$ . It follows from the above that to any  $y \in \text{Ob}(\mathcal{D}')$  we can associate some (uniquely defined) object to be denoted  $Q(y)$ . We now should define  $Q$  on morphisms. Consider a map  $y \rightarrow z$  for some  $y, z \in \text{Ob}(\mathcal{D}')$ . We have already realised  $Q(y), Q(z)$  as coequalizers of certain diagrams. It is clear that  $f$  induces a morphism of these diagrams. So we obtain a desired morphism  $Q(f)$ .  $\square$

Let us now apply lemma 3.60 to  $R' = \text{Forg}_{\mathcal{C}}, L' = \tilde{L}$ . We obtain a functor  $Q: T - \text{mod}_{\mathcal{C}} \rightarrow \mathcal{D}$ . Note now that the functor  $\tilde{R}: \mathcal{D} \rightarrow T - \text{mod}_{\mathcal{C}}$  coincides with functor from lemma 3.60 for a pair  $(\text{Forg}_{\mathcal{C}}, \tilde{L}), (R, L)$  (we can apply this lemma to this pair because it follows from the proof of the implication  $\Rightarrow$  that the condition (2) of theorem 3.59 holds for  $\text{Forg}_{\mathcal{C}}$ ). It now follows from the uniqueness part of lemma 3.60 that  $Q \circ \tilde{R} = id_{\mathcal{D}}, \tilde{R} \circ Q = id_{T - \text{mod}_{\mathcal{C}}}$ .  $\square$

### Remark 3.61

Note that if  $\mathcal{D}$  is (finite) cocomplete and  $\mathcal{F}$  is cocontinuous then (2) holds automatically.

The following lemma is very useful for checking condition (1) of Theorem 3.59.

### Lemma 3.62

Suppose  $\mathcal{D}$  is abelian and  $R$  is right exact. Then  $R$  is conservative iff for any  $X$  with  $R(X) \simeq 0$  we have  $X \simeq 0$ .

*Proof.* Implication  $\Rightarrow$  is obvious. Let us prove the implication  $\Leftarrow$ . Suppose that we have  $f: X \rightarrow Y$  such that  $R(f)$  is an isomorphism. Note that  $R$  being right exact and right adjoint is exact. It follows that  $R(\ker f) = \ker R(f) = 0, R(\text{coker } f) = \text{coker } R(f) = 0$  so  $\ker f \simeq \text{coker } f \simeq 0$ . We conclude that  $f$  is an isomorphism.  $\square$

We are now answering to the question when  $T - \text{mod}_{\mathcal{A}}$  is abelian.

### Proposition 3.63

If  $\mathcal{A}$  is an abelian category and  $T: \mathcal{A} \rightarrow \mathcal{A}$  is a right exact monad on  $\mathcal{A}$  then  $\mathcal{B} := T - \text{mod}_{\mathcal{A}}$  is abelian.

*Proof.* Let us first of all show that the category  $T - \text{mod}_{\mathcal{A}}$  is pre-abelian.

To do this we need to construct kernels and co-kernels in the category  $T - \text{mod}_{\mathcal{A}}$ . If  $h: X \rightarrow X'$  induces a morphism  $(X, f) \rightarrow (X', f')$  in  $\mathcal{D}$  then the kernel of this morphism is exactly  $(\ker h, l)$ , where  $l: T(\ker h) \rightarrow \ker h$  is obtained as the composition of the morphisms  $T(\ker h) \rightarrow \ker T(h) \rightarrow \ker h$ . To construct co-kernel we recall that  $T$  is right exact so  $T(\text{coker } h) = \text{coker } T(h)$  and we can define cokernel of the morphism  $h: (X, f) \rightarrow (X', f')$  to be  $(\text{coker}(h), p)$ , where  $p: T(\text{coker } h) = \text{coker } T(h) \rightarrow \text{coker}(h)$  is the natural morphism.

Consider now a forgetful functor  $\text{forg}: \mathcal{B} \rightarrow \mathcal{A}$ . It follows from the constructions that  $\text{forg}$  is an exact functor. Note that  $\text{forg}$  is conservative and  $\mathcal{C}$  is abelian so it follows that the comparison morphism

$$\text{coker}(\ker(h)) \rightarrow \ker(\text{coker}(h))$$

is an isomorphism in  $\mathcal{D}$ , hence,  $\mathcal{D}$  is abelian. □

#### 4. DELIGNE-KELLY TENSOR PRODUCT

Recall that on the previous lecture factorization homologies were constructed. They were constructed as a colimit along some diagram. So we want to work with categories which have enough colimits (cocomplete). So a good setting for us will be  $\mathbf{Pr}_{\mathcal{C}}$ . Recall that by proposition 2.43 we have an equivalence  $\mathbf{Pr}_{\mathcal{C}} \simeq \mathbf{Rex}$ . We want to have a tensor structure on the category  $\mathbf{Rex} \simeq \mathbf{Pr}_{\mathcal{C}}$ .

##### Definition 4.64

Given two  $k$ -linear categories  $\mathcal{A}, \mathcal{B}$  we define their tensor product  $\mathcal{A} \otimes \mathcal{B}$  as a category with objects the pairs  $(A, B)$  with  $A \in \text{Ob}(\mathcal{A}), B \in \text{Ob}(\mathcal{B})$  and morphisms  $\text{Hom}((A, B), (A', B')) := \text{Hom}(A, B) \otimes_k \text{Hom}(A', B')$ .

##### Proposition 4.65

(a) Category  $\mathcal{A} \otimes \mathcal{B}$  has the following universal property – we have a natural equivalence between  $k$ -bilinear functors  $\mathcal{A} \times \mathcal{B} \rightarrow \mathcal{C}$  to some  $k$ -linear category  $\mathcal{C}$  and functors  $\mathcal{A} \otimes \mathcal{B} \rightarrow \mathcal{C}$ .

(b) For any  $k$ -linear  $\mathcal{A}, \mathcal{B}, \mathcal{C}$  we have  $[\mathcal{A} \otimes \mathcal{B}, \mathcal{C}] \simeq [\mathcal{A}, [\mathcal{B}, \mathcal{C}]]$ .

It turns out that if  $\mathcal{A}, \mathcal{B}$  are cocomplete then the category  $\mathcal{A} \otimes \mathcal{B}$  need not to be cocomplete. There is a way to define other tensor product to be denoted by  $\boxtimes$  such that tensor product of two cocomplete categories will be cocomplete. References are [K, Section 6.5], [S, Section 2.3], [EGNO, Section 1.11],

##### Theorem 4.66

(a) For any two  $\mathcal{A}, \mathcal{B} \in \mathbf{Rex}$  there exists a category  $\mathcal{A} \boxtimes \mathcal{B}$  uniquely defined by the following property

$$\mathbf{Rex}[\mathcal{C} \boxtimes \mathcal{D}, \mathcal{E}] \simeq \text{Bilin}(\mathcal{C} \times \mathcal{D}, \mathcal{E}),$$

where  $\text{Bilin}(\mathcal{C} \times \mathcal{D}, \mathcal{E})$  is the category of  $k$ -bilinear functors preserving finite colimits in each variable.

We have an equivalence

$$\mathbf{Rex}[\mathcal{C} \boxtimes \mathcal{D}, \mathcal{E}] \simeq \mathbf{Rex}[\mathcal{C}, \mathbf{Rex}[\mathcal{D}, \mathcal{E}]].$$

(b) The Kelly tensor product  $\boxtimes$  equips  $\mathbf{Rex}$  with the structure of a symmetric closed monoidal  $(2, 1)$ -category.

(c) The tensor product  $\boxtimes$  extends to a monoidal structure on  $\mathbf{Pr}_c$  and the functor  $\text{ind}$  extends to an equivalence  $\mathbf{Rex}^{\boxtimes} \xrightarrow{\sim} \mathbf{Pr}_c^{\boxtimes}$  of symmetric monoidal  $(2, 1)$ -categories.

**Remark 4.67**

Assume that  $\mathcal{C}, \mathcal{D}$  are locally finite abelian. Then by proposition 3.56 we have an equivalence  $\mathcal{C} \simeq C - \text{comod}$ ,  $\mathcal{D} \simeq D - \text{comod}_{f.d.}$  for some coalgebras  $C, D$  over  $k$ . Then we have  $\mathcal{C} \boxtimes \mathcal{D} \simeq (C \otimes D) - \text{comod}$  so we have a rather explicit description of the category  $\mathcal{C} \boxtimes \mathcal{D}$  in this case.

## 5. MONADICITY FOR MODULE CATEGORIES

Main reference for all the notions of this section is [EGNO, Chapter 7]. The main theorem is taken from [BBJ, Theorem 4.6].

**Definition 5.68**

Tensor category  $(\mathcal{C}, \otimes)$  is called closed if for any  $X \in \text{Ob}(\mathcal{C})$  the functor  $X \otimes \bullet: \mathcal{C} \rightarrow \mathcal{C}$  admits a right adjoint. If  $\mathcal{C} \in \mathbf{Pr}$  then this is equivalent to the fact that  $X \otimes \bullet$  is cocomplete.

**Definition 5.69**

An object  $C \in \text{Ob}(\mathcal{C})$  of a tensor category is called right dualizable if there exists an object  $C^* \in \text{Ob}(\mathcal{C})$  and morphisms  $\text{ev}_C: C^* \otimes C \rightarrow 1$ ,  $\text{coev}_C: 1 \rightarrow C \otimes C^*$  such that the compositions

$$\begin{aligned} C &\xrightarrow{\text{coev}_C \otimes \text{id}_C} (C \otimes C^*) \otimes C \xrightarrow{\alpha_{C, C^*, C}} C \otimes (C^* \otimes C) \xrightarrow{\text{id}_C \otimes \text{ev}_C} C, \\ C^* &\xrightarrow{\text{id}_{C^*} \otimes \text{coev}_C} C^* \otimes (C \otimes C^*) \xrightarrow{\alpha_{C^*, C, C^*}^{-1}} (C^* \otimes C) \otimes C \xrightarrow{\text{ev}_C \otimes \text{id}_C} C \end{aligned}$$

are identity morphisms, here  $\alpha$  are associators for the tensor structure  $\otimes$  on  $\mathcal{C}$ .

An object  $C \in \text{Ob}(\mathcal{C})$  of a tensor category is called left dualizable if there exists an object  $*C \in \text{Ob}(\mathcal{C})$  and morphisms  $\text{ev}'_C: C \otimes *C \rightarrow 1$ ,  $\text{coev}'_C: 1 \rightarrow *C \otimes C$  such that the compositions

$$\begin{aligned} C &\xrightarrow{\text{id}_C \otimes \text{coev}'_C} C \otimes (*C \otimes C) \xrightarrow{\alpha_{C, *C, C}} (C \otimes C^*) \otimes C \xrightarrow{\text{id}_C \otimes \text{ev}_C} C, \\ *C &\xrightarrow{\text{coev}'_C \otimes \text{id}_{C^*}} (*C \otimes C) \otimes *C \xrightarrow{\alpha_{C^*, C, C^*}^{-1}} *C \otimes (C \otimes *C) \xrightarrow{\text{id}_C \otimes \text{ev}'_C} *C. \end{aligned}$$

**Definition 5.70**

Tensor category  $(\mathcal{C}, \otimes)$  is called rigid if all compact objects of  $\mathcal{C}$  are right and left dualizable.

**Definition 5.71**

Let  $(\mathcal{A}, \otimes)$  be a tensor category in a category  $\mathbf{Pr}_c$ .

(1) A (right)  $\mathcal{A}$ -module category  $\mathcal{M}$  in  $\mathbf{Pr}_c$  is a category  $\mathcal{M} \in \mathbf{Pr}_c$  together with an action functor

$$\text{act}_{\mathcal{M}}: \mathcal{M} \boxtimes \mathcal{A} \rightarrow \mathcal{M}$$

satisfying standard associativity (pentagon) axioms (this notion categorifies the notion of a module over an algebra). We will denote  $act_{\mathcal{M}}(m \boxtimes X)$  by  $m \otimes X$ ,  $m \in \text{Ob}(\mathcal{M})$ ,  $X \in \text{Ob}(\mathcal{C})$ .

(2) For any  $m \in \mathcal{M}$  the functor  $act_m: \mathcal{A} \rightarrow \mathcal{M}$ ,  $a \mapsto m \otimes a$  admits a right adjoint to be denoted  $act_m^R: \mathcal{M} \rightarrow \mathcal{A}$  (by proposition 1.32).

(3) For  $m, n \in \mathcal{M}$  we set  $\underline{\text{Hom}}_{\mathcal{A}}(m, n) := act_m^R(n) \in \mathcal{A}$ . Note that the object  $\underline{\text{Hom}}(m, n)$  (contravariantly) represents the functor  $\mathcal{A} \rightarrow \text{Vect}_k$ ,  $X \mapsto \text{Hom}_{\mathcal{M}}(X \otimes m, n)$  i.e.

$$\text{Hom}_{\mathcal{A}}(X, \underline{\text{Hom}}(m, n)) = \text{Hom}_{\mathcal{A}}(X, act_m^R(n)) \simeq \text{Hom}_{\mathcal{M}}(m \otimes X, n). \quad (5.1)$$

The object  $\underline{\text{Hom}}(m, n)$  is called internal  $\text{Hom}$  form  $m$  to  $n$ .

(4) For any triple  $m, n, p \in \mathcal{M}$  there is a well-defined composition map in  $\mathcal{A}$ :

$$\underline{\text{Hom}}(n, p) \otimes \underline{\text{Hom}}(m, n) \rightarrow \underline{\text{Hom}}(m, p). \quad (5.2)$$

To construct it let us note that for any  $m, n \in \text{Ob}(\mathcal{M})$  we have a canonical evaluation morphism

$$ev_{m,n}: \underline{\text{Hom}}_{\mathcal{A}}(m, n) \otimes m \rightarrow n$$

which corresponds to  $id_{\underline{\text{Hom}}_{\mathcal{A}}(m, n)}$  under the isomorphism (5.1). We can now form the composition

$$\begin{aligned} (\underline{\text{Hom}}_{\mathcal{A}}(n, p) \otimes \underline{\text{Hom}}_{\mathcal{A}}(m, n)) \otimes m &\xrightarrow{\alpha_{\underline{\text{Hom}}_{\mathcal{A}}(n, p), \underline{\text{Hom}}_{\mathcal{A}}(m, n), m}} \\ &\underline{\text{Hom}}_{\mathcal{A}}(n, p) \otimes (\underline{\text{Hom}}_{\mathcal{A}}(m, n) \otimes m) \xrightarrow{id \otimes ev_{m, n}} \\ &\underline{\text{Hom}}_{\mathcal{A}}(n, p) \otimes n \xrightarrow{ev_{n, p}} p \end{aligned}$$

which gives us the composition map (5.2) via the isomorphism (5.1).

(5) We set  $\underline{\text{End}}(m) := \underline{\text{Hom}}(m, m) = act_m^R(m) = act_m^R(act_m(1))$ . By (4),  $\underline{\text{End}}(m)$  carries a unital algebra structure i.e.  $\underline{\text{End}}(m) \in \mathcal{A}$  is a unital algebra object. Recall that by  $\underline{\text{End}}(m) - \text{mod}_{\mathcal{A}}$  we denote the category of  $\underline{\text{End}}(m)$ -modules (in  $\mathcal{A}$ ).

(6) We say that  $m$  is an  $\mathcal{A}$ -generator if  $act_m^R$  is faithful.

(7) We say that  $m$  is an  $\mathcal{A}$ -projective if  $act_m^R$  is colimit-preserving (this is equivalent to say that  $act_m^R$  preserves finite colimits since  $act_m^R$  preserves filtered colimits and now apply remark 1.17).

(8) We say that  $m$  is an  $\mathcal{A}$ -progenerator if it is an  $\mathcal{A}$ -projective  $\mathcal{A}$ -generator.

### Definition 5.72

Let  $\mathcal{M}, \mathcal{N}$  be  $\mathcal{A}$ -module categories. A  $\mathcal{A}$ -linear functor from  $\mathcal{M}$  to  $\mathcal{N}$  consists of a functor  $\mathcal{F}: \mathcal{M} \rightarrow \mathcal{N}$  and a natural isomorphism  $s_{X, M}: \mathcal{F}(X \otimes M) \xrightarrow{\sim} X \otimes \mathcal{F}(M)$ ,  $X \in \text{Ob}(\mathcal{A})$ ,  $M \in \text{Ob}(\mathcal{M})$  which is associative and compatible with tensor product by  $1 \in \text{Ob}(\mathcal{C})$ :

### Theorem 5.73

(Monadicity for module categories) Let  $\mathcal{A}$  be a rigid abelian tensor category in  $\mathbf{Pr}_c$  and let  $\mathcal{M} \in \mathbf{Pr}_c$  be an abelian  $\mathcal{A}$ -module category with an  $\mathcal{A}$ -progenerator  $m \in \mathcal{M}$ .

Then the functor  $\widetilde{\text{act}}_m^R(\underline{\text{Hom}}(m, \bullet))$  is an equivalence of  $\mathcal{A}$ -module categories,

$$\mathcal{M} \simeq \underline{\text{End}}(m) - \text{mod}_{\mathcal{A}},$$

where  $\mathcal{A}$  acts on the right by multiplication.

*Proof.* (Sketch) We want to apply theorem 3.59 to the functor  $\text{act}_m: \mathcal{A} \rightarrow \mathcal{M}$ . Recall that  $m$  is an  $\mathcal{A}$ -generator so by lemma 3.62 the functor  $\text{act}_m^R$  is conservative. Note also that  $m$  is  $\mathcal{A}$ -projective i.e. it is colimit-preserving in particular it preserves coequalizers. So by theorem 3.59 we obtain an equivalence

$$\mathcal{M} \xrightarrow{\sim} \text{act}_m^R \circ \text{act}_m - \text{mod}_{\mathcal{A}}.$$

It remains to identify  $\text{act}_m^R \circ \text{act}_m - \text{mod}_{\mathcal{A}} \simeq \underline{\text{End}}(m) - \text{mod}_{\mathcal{A}}$ , functor  $\widetilde{\text{act}}_m^R$  will identify with  $\underline{\text{Hom}}(m, \bullet)$ .

One can show that  $\text{act}_m^R: \mathcal{M} \rightarrow \mathcal{A}$  carries a canonical module structure so the composition  $\text{act}_m^R \circ \text{act}_m: \mathcal{A} \rightarrow \mathcal{A}$  is a module functor. Any module functor  $\mathcal{F}: \mathcal{A} \rightarrow \mathcal{A}$  is isomorphic to the functor  $\mathcal{F}(1) \otimes \bullet$  (this is a categorification of the fact that if  $A$  is a  $k$ -algebra then any (right)  $A$ -module homomorphism  $f: A \rightarrow A$  is given by  $a \mapsto f(1)a$ ). So we see that  $\text{act}_m^R \circ \text{act}_m \simeq \text{act}_m^R \circ \text{act}_m(1) \otimes \bullet = \underline{\text{End}}(m) \otimes \bullet$  and the claim follows.

One can also show directly (in the spirit of proposition 3.53) that the functor  $\underline{\text{Hom}}(m, \bullet)$  defines an equivalence  $\mathcal{M} \xrightarrow{\sim} \underline{\text{End}}(m) - \text{mod}_{\mathcal{A}}$  (see [EGNO, section 7.10]).  $\square$

#### Remark 5.74

Recall that  $\underline{\text{End}}(m) - \text{mod}_{\mathcal{A}}$  is the category of  $\underline{\text{End}}(m)$ -modules in the tensor category  $\mathcal{A}$ .

**Example 5.75.** One can hold in head the following example. Let  $\mathcal{A} = \text{Vect}_{k,f.d.}$  be the category of finite dimensional vector spaces over  $k$  and  $\mathcal{M}$  is some finite category. We consider  $\mathcal{A}, \mathcal{M}$  as objects of the tensor category  $(\mathbf{Cat}, \times)$  (analogy of  $\mathbf{Pr}_c$ ). Then the category  $(\text{Vect}_{k,f.d.}, \otimes)$  is a tensor category in  $\mathbf{Cat}$  and the natural functor  $\text{Vect}_{k,f.d.} \times \mathcal{M} \rightarrow \mathcal{M}$  (see remark 1.2) equips  $\mathcal{M}$  with a structure of  $\text{Vect}_{k,f.d.}$ -module. For  $M \in \mathcal{M}$  the functor  $\text{act}_M: \text{Vect}_{k,f.d.}$  sends  $V$  to  $V \otimes_k M$  the right adjoint  $\text{act}_M^R: \mathcal{M} \rightarrow \text{Vect}_{k,f.d.}$  sends  $N$  to  $\text{Hom}_{\mathcal{M}}(M, N) \in \text{Ob}(\text{Vect}_{k,f.d.})$ .

Now the theorem exactly says that if  $P \in \mathcal{M}$  is a projective generator then  $\mathcal{M} \simeq \underline{\text{End}}(P) - \text{mod}_{\text{Vect}_{f.g.}} = \text{End}(P) - \text{mod}_{f.g.}$  (c.f. proposition 3.53).

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